

Introduction

- In reservoir planning and operation problem, inflow, Q_p is a random variable and is not known with certainty. Its probability distribution, however, may be estimated from the historical sequence of inflows.
- Being functions of inflow, Q_t , storage S_t and release R_t are also random variables.
- In a constraint containing two random variables, if the probability distribution of one is known, the probabilistic behavior of the second can be expressed as a measure of chance in terms of the probability of the first variable.
- If a constraint contains more than two random variables, we get into computational complications, and we need to understand the specific problem clearly to reformulate the problem, if necessary, and avoid those complications.

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Chance Constraint

- The constraint, relating the release, R_p (random) and demand, D_p (deterministic), is expressed as a chance constraint, $P(R_i \ge 0) \ge \alpha_i$: It means that the probability of release equaling or exceeding the known demand is at least equal to α_1 , which is referred to as the reliability level. The interpretation of this chance constraint is simply that the reliability of meeting the demand in peniod t is at
- least α_i . Similarly, Chance constraints for the maximum release and the maximum and minimum storage can be written as $P[R_i \le R_i^{max}] \ge \alpha_2$

- $F(K_1, S_1, K_2) = 2G_2$ $F(S_1, S_1, K_2) \ge G_2$ and $F(S_2, S_2) \ge G_2$ To use the above chance constraints in an optimization algorithm, we must first determine the probability distribution of R_1 and S_1 from the known probability distribution of Q_2 . However, because S_1 , G_2 , and G_3 are all interdependent through the continuity equation, it is, in general, not possible to derive the probability distributions of both S_3 and G_3 simultaneously.
- Stritulaneously.

 To overcome this difficulty and to enable the use of linear programming in the solution, linear decision rule is appropriately defined.

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Linear Decision Rule (LDR)

The linear decision rule (LDR) relates the release, $R_{\rm p}$ from the reservoir as a linear function of the water available in period t. The simplest form of such an LDR is

where b_t is a deterministic parameter called the decision parameter.

In this LDR, the entire amount, $Q_{\rm e}$ is taken into account while making the release decision. Depending on the proportion of inflow, $Q_{\rm e}$ used in the linear decision rule, a number of such LDRs may be formulated. A general form of this LDR may be written as

 $R_t = S_t + \beta_t Q_t - b_t$ $0 \le \beta_t \le 1$

 β_i = 0 yields a relatively conservative release policy with release decisions related only to the storage, S_i β_i = 1 yields an optimistic policy where the entire amount of water available (S_i+Q_i) is used in the LDR.

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LDR - Contd.

 $R_{i} = S_{i} + Q_{i} - b_{i}$

Storage continuity equation is $S_{t+1} = S_t + Q_t - R_t$ Using above two equations

- $S_{t+1} = b_t$ Thus, the random variable, S_{t+1} , is set equal to a deterministic parameter b_r Thus, the role of the LDR in this case is to treat S_r deterministic in formulation. A main advantage of doing this is to do away with one of the random variables, S_p so that the distribution of the other random variable, R_r may be expressed in terms of the known distribution of Q_r .
- This implies that the variance of Q_t is entirely transferred to the variance of R_t
- Including evaporation loss as a storage-dependent term in the storage continuity equation, the linear decision rule is written as,

$$R_t = Q_t - [A_0 e_t] + [1 - (a e_t/2)] b_{t-1} - [1 + (a e_t/2)] b_t$$

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Deterministic Equivalent of a Chance Constraint

Knowing the probability distribution of inflow, $Q_{\rm I}$, it is possible to obtain the deterministic equivalents of the chance constraints using the LDR, as follows:

$$\begin{split} &P[R_t \geq D_t] \geq \alpha_1 \\ &P[S_t + Q_t - b_t \geq D_t] \geq \alpha_1 \end{split}$$
 $P[b_{t-1} + Q_t - b_t \ge D_t] \ge \alpha_1$ $P[Q_t \ge D_t + b_t - b_{t-1}] \ge \alpha_1$ $P[Q_t \le D_t + b_t - b_{t-1}] \le 1 - \alpha_1$



The term $D_l + b_l - b_{l-1}$, is deterministic with b_l and b_{l-1} being decision variables and D_l being a known quantity for the period t.

$$F_{Qt}(D_t + b_t - b_{t-1}) \le 1 - \alpha_1$$

 $(D_t + b_t - b_{t-1}) \le F_{Qt}^{-1}(1 - \alpha_1)$

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